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CITATION:

Miyakawa, Kampei ...[et al]. On Extremal Problems of MPR-posets II (Models of Computation and Algorithms). 数理解析研究所講究録 1999, 1093: 57-61

ISSUE DATE:

1999-04

URL:

<http://hdl.handle.net/2433/62972>

RIGHT:

## On Extremal Problems of MPR-posets II

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A mathematical theory for the subject on ancestral character-state reconstructions under the maximum parsimony in phylogeny has been developing [1-9]. In this paper, we show some extremal properties of  $\sigma(r)$ -version MPR-posets, particularly, the lattice-theoretic properties of those posets.

We use the notations in [1, 4]. Let  $T = (V, E, \sigma)$  be any undirected tree whose endnodes are evaluated by a weight function  $\sigma : V_O \rightarrow \Omega$ , where  $\Omega$  expresses the linearly ordered character-states. From the viewpoint of enumeration, let  $\Omega$  denote the set of non-negative integers.  $V$  is the set of nodes,  $V_O$  is the set of endnodes which are nodes of degree one,  $V_H$  is the set of internal nodes, and  $E$  is the set of branches. Note that  $V_O \cup V_H = V$  and  $V_O \cap V_H = \emptyset$ . We call this tree an *el-tree*. For an el-tree  $T$ , we define an assignment  $\lambda : V \rightarrow \Omega$  such that  $\lambda|_{V_O}$  (the restriction of  $\lambda$  to  $V_O$ ) =  $\sigma$ , where  $\lambda(u)$  is called a *state* of  $u$  under  $\lambda$ . This assignment is called a *reconstruction* on an el-tree  $T$ . For each branch  $e$  in  $E$  of an el-tree  $T$  with a reconstruction  $\lambda$ , we define the *length*  $l(e)$  of branch  $e = \{u, v\}$  by  $|\lambda(u) - \lambda(v)|$ . Then the length  $L(T|\lambda)$  of an el-tree  $T$  under the reconstruction  $\lambda$  is the sum of the lengths of the branches. That is,  $L(T|\lambda) = \sum_{e \in E} l(e)$ . Furthermore, we define the minimum length  $L^*(T)$  of  $T$  by

$$L^*(T) = \min\{L(T|\lambda) \mid \lambda \text{ is a reconstruction on } T\}.$$

Note that  $L^*(T)$  is well-defined. A reconstruction  $\lambda$  such that  $L(T|\lambda) = L^*(T)$  is called a *most-parsimonious reconstruction* (abbreviated to MPR) on  $T$ . Note that generally an el-tree has more than one MPR. We denote the set of MPRs on  $T$  by  $\mathbf{Rmp}(T)$ . For each node  $u$  in  $V$ , the set  $\{\lambda(u) \mid \lambda \in \mathbf{Rmp}(T)\}$  of states of  $u$  under  $\lambda$  such that  $\lambda \in \mathbf{Rmp}(T)$  is called the *MPR-set* of  $u$  and written as  $S_u$ . The algorithms to get  $L^*(T)$ ,  $\mathbf{Rmp}(T)$ , and  $S_u$  are given in [1, 4]. See [1, 4] for details.

For a given el-tree  $T = (V, E, \sigma(r))$ , we define a *rooted el-tree*  $T^{(r)}$  rooted at any element  $r$  in  $V$ . If  $r$  is an endnode, i.e.,  $r \in V_O$  and  $s$  is its unique child, we denote the rooted el-tree  $T^{(r)}$  by  $(T_s, r)$ . The parent-child relation  $\{u, v\}$  in  $E$  in a rooted el-tree  $T$  is denoted by  $u \rightarrow v$  which means  $u$  is a parent of  $v$ . Let  $I_i$  ( $i \in A$ ) be any family of closed intervals in  $\Omega$ . Let the median two points of all the endpoints of  $I_i$  be  $\langle x, y \rangle$ . Then we define the *median interval* of  $I_i$  ( $i \in A$ ) by the closed interval  $[x, y]$  and denote by  $\text{med}\langle I_i : i \in A \rangle$ . For each node  $u$  in the body of a rooted el-tree  $T$ , we assign a closed interval  $I(u)$  of  $\Omega$  recursively as follows:

$$I(u) = \begin{cases} [\sigma(u), \sigma(u)] & \text{if } u \text{ is a leaf,} \\ \text{med}\langle I(v) : u \rightarrow v \rangle & \text{otherwise.} \end{cases}$$

This interval  $I(u)$  is called the *characteristic interval* of a node  $u$ .

From a phylogenetic point of view, Minaka [2,3] has introduced the two partial orderings on  $\mathbf{Rmp}(T)$  to investigate the relationships among MPRs. One is the usual ordering, and the other is a partial ordering that depends on a state of a specified root of a given el-tree. We now give a mathematically explicit formulation for those partial orderings. Let  $T$  be an el-tree. The usual ordering  $\lambda \leq \mu$  on  $\mathbf{Rmp}(T)$  is defined by  $\lambda(u) \leq \mu(u)$  for all  $u$  in  $V$ . Let  $T$  be a rooted el-tree  $(T_s, r)$ . A binary relation  $a \leq_{\sigma(r)} b$  on  $\Omega$  is defined by  $\sigma(r) \leq a \leq b$  or  $\sigma(r) \geq a \geq b$ . Then, a binary relation  $\lambda \leq_{\sigma(r)} \mu$  on  $\mathbf{Rmp}(T)$  is defined by  $\lambda(u) \leq_{\sigma(r)} \mu(u)$  for all  $u$  in  $V$ . It is easily shown that those relations are partial orderings. The partially ordered set  $(\mathbf{Rmp}(T), \leq)$  is called a *usual MPR-poset*, and the  $(\mathbf{Rmp}(T), \leq_{\sigma(r)})$  is called a  $\sigma(r)$ -version MPR-poset. Note that a usual MPR-poset is uniquely defined for an el-tree, but a  $\sigma(r)$ -version MPR-poset, depending on a specified root, is defined in several ways for an el-tree.

We first restate some previous results which relate particularly to new results stated later. Let  $T$  be a rooted el-tree  $(T_s, r)$ . We define a reconstruction  $\lambda$  on  $T$  by  $\lambda(u) = x$  in  $S_u$  satisfying  $x \leq_{\sigma(r)} y$  for any  $y$  in  $S_u$ , that is,  $x$  is the least element of a subposet  $(S_u, \leq_{\sigma(r)})$  in the poset  $(\Omega, \leq_{\sigma(r)})$ . This reconstruction  $\lambda$  is well-defined since it is easily seen that the least element of each subposet  $(S_u, \leq_{\sigma(r)})$ , and then this  $\lambda$  is particularly written as  $\lambda_{\min}^{<\sigma(r)>}$ . The following theorem answers for whether there exists the least element in a  $\sigma(r)$ -version MPR-poset or not.

**Theorem A.** *Let  $T$  be a rooted el-tree  $(T_s, r)$ . Then the reconstruction  $\lambda_{\min}^{<\sigma(r)>}$  is the least element of  $(\mathbf{Rmp}(T), \leq_{\sigma(r)})$ .  $\square$*

It is known that  $(\mathbf{Rmp}(T), \leq_{\sigma(r)})$  doesn't always have the greatest element. The following shows one of the requirements for a reconstruction  $\lambda$  in  $\mathbf{Rmp}(T)$  to be a maximal element of the  $\sigma(r)$ -version MPR-poset. Let  $T$  be a rooted el-tree  $(T_s, r)$ . We define two reconstructions  $\alpha^{<\sigma(r)>}$  and  $\beta^{<\sigma(r)>}$  on  $T$  by  $\alpha^{<\sigma(r)>}(u) =$  the smallest element, under the usual ordering  $\leq$ , of maximal elements in the subposet  $(S_u, \leq_{\sigma(r)})$  and  $\beta^{<\sigma(r)>}(u) =$  the greatest element, under the usual ordering  $\leq$ , of maximal elements in the subposet  $(S_u, \leq_{\sigma(r)})$ .

**Proposition B.** *Let  $T$  be a rooted el-tree  $(T_s, r)$ . Then, both  $\alpha^{<\sigma(r)>}$  and  $\beta^{<\sigma(r)>}$  are maximal elements of  $(\mathbf{Rmp}(T), \leq_{\sigma(r)})$ .  $\square$*

The following shows a necessary and sufficient condition for a  $\sigma(r)$ -version MPR-poset to have the greatest element.

**Corollary C.** *Let  $T$  be a rooted el-tree  $(T_s, r)$ .  $(\mathbf{Rmp}(T), \leq_{\sigma(r)})$  has the greatest element if and only if for any  $u$  in  $V_H$ ,  $\sigma(r) \leq \min(S_u)$  or  $\sigma(r) \geq \max(S_u)$ .  $\square$*

Using those results, we have some new results about the characteristics of MPR-posets. The following theorem shows the necessary and sufficient condition for an MPR to be a

maximal element of  $\sigma(r)$ -version MPR-poset.

**Theorem 1.** *Let  $T$  be a rooted el-tree  $(T_s, r)$  and  $\lambda$  in  $\mathbf{Rmp}(T)$ .  $\lambda$  is a maximal element of  $(\mathbf{Rmp}(T), \leq_{\sigma(r)})$  if and only if for any  $u$  in  $V_H$ ,  $\lambda(u)$  is a maximal element of subposet  $(S_u, \leq_{\sigma(r)})$ .  $\square$*

Note that there exists an el-tree  $T$  such that the number of all maximal elements of  $\sigma(r)$ -version MPR-poset is exponential for the number  $n$  of the nodes. For example, the rooted el-tree  $T = (T_s, r)$  shown in Fig.1 has  $6m+3$  nodes, and the  $\sigma(r)$ -version MPR-poset  $(\mathbf{Rmp}(T), \leq_{\sigma(r)})$  has  $2^m + 1$  maximal elements.

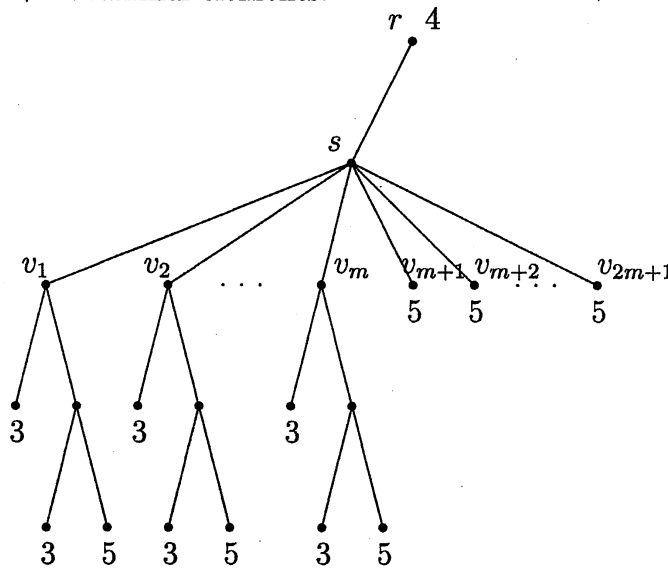


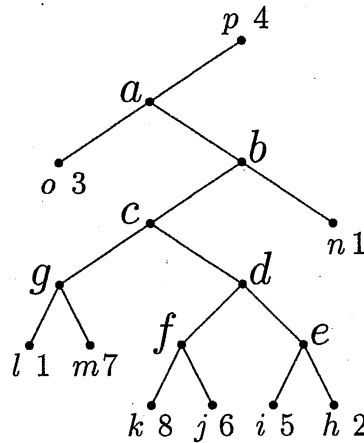
Fig. 1: A rooted el-tree with  $6m + 3$  nodes

We next show the followings which answer whether any  $\sigma(r)$ -version MPR-poset forms a lattice or not.

**Theorem 2.** *Let  $T$  be a rooted el-tree  $(T_s, r)$ . The  $\sigma(r)$ -version MPR-poset  $(\mathbf{Rmp}(T), \leq_{\sigma(r)})$  forms a lower semilattice.  $\square$*

**Theorem 3.** *Let  $T$  be a rooted el-tree  $(T_s, r)$ . The  $\sigma(r)$ -version MPR-poset  $(\mathbf{Rmp}(T), \leq_{\sigma(r)})$  forms an upper semilattice if and only if for any  $u$  in  $V_H$ ,  $\sigma(r) \leq \min(S_u)$  or  $\max(S_u) \leq \sigma(r)$  holds.  $\square$*

We here show some examples of the theorems stated above. Let  $(T_a, p)$  be an el-tree  $T$  rooted at  $p$  shown in Fig.2. From  $\mathbf{Rmp}(T)$  shown in Table 1, we can construct a  $\sigma(p)$ -version MPR-poset  $(\mathbf{Rmp}(T), \leq_{\sigma(p)})$  shown in Fig.3, whose maximal elements  $\lambda_1$ ,  $\lambda_3$ , and  $\lambda_6$  assign a maximal element of  $(S_u, \leq_{\sigma(p)})$  for each node  $u$ , and the poset forms a lower semilattice.

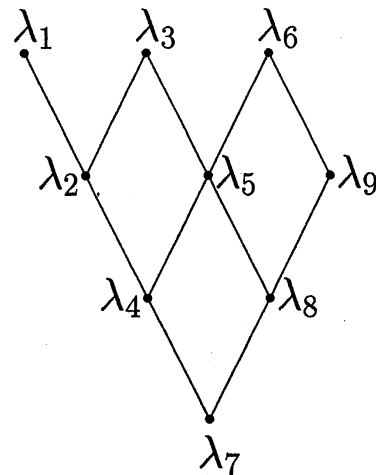
Fig. 2: a rooted el-tree  $(T_a, p)$ 

We finally show the following, which is immediate from Corollary C, Theorem 2, and Theorem 3.

**Corollary 1.** *Let  $T$  be a rooted el-tree  $(T_s, r)$ . Then the following three statements are equivalent:*

1. *The  $\sigma(r)$ -version MPR-poset has the greatest element.*
2. *The  $\sigma(r)$ -version MPR-poset forms an upper-semilattice.*
3. *The  $\sigma(r)$ -version MPR-poset forms a lattice.*  $\square$

$\lambda$	a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p
$\lambda_1$	3	3	3	3	3	6	3	2	5	6	8	1	7	1	3	4
$\lambda_2$	3	3	3	4	4	6	3	2	5	6	8	1	7	1	3	4
$\lambda_3$	3	3	3	5	5	6	3	2	5	6	8	1	7	1	3	4
$\lambda_4$	3	3	4	4	4	6	4	2	5	6	8	1	7	1	3	4
$\lambda_5$	3	3	4	5	5	6	4	2	5	6	8	1	7	1	3	4
$\lambda_6$	3	3	5	5	5	6	5	2	5	6	8	1	7	1	3	4
$\lambda_7$	4	4	4	4	4	6	4	2	5	6	8	1	7	1	3	4
$\lambda_8$	4	4	4	5	5	6	4	2	5	6	8	1	7	1	3	4
$\lambda_9$	4	4	5	5	5	6	5	2	5	6	8	1	7	1	3	4

Table 1:  $\mathbf{Rmp}(T)$ Fig. 3:  $(\mathbf{Rmp}(T), \leq_{\sigma(p)})$

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